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(54) **METHOD OF MALWARE DETECTION AND SYSTEM THEREOF**

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This patent is subject to a terminal disclaimer.

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Installing a Tanium Client—distributed also using user/password (not encrypted) from a hosted Module server, downloaded from https://docs.tanium.com/client/client/deployment.html#client_management on Apr. 1, 2021.

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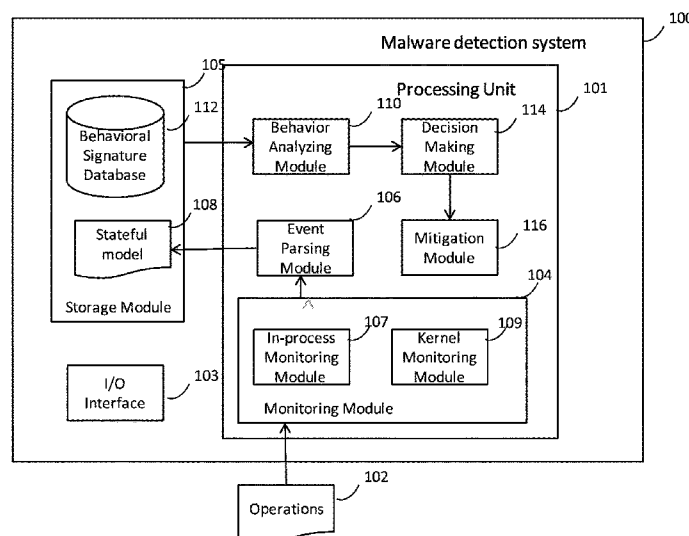
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(57) **ABSTRACT**

There is provided a system and a computer-implemented method of detecting malware in real time in a live environment. The method comprises: monitoring one or more operations of at least one program concurrently running in the live environment, building at least one stateful model in accordance with the one or more operations, analyzing the at least one stateful model to identify one or more behaviors, and determining the presence of malware based on the identified one or more behaviors.

20 Claims, 5 Drawing Sheets



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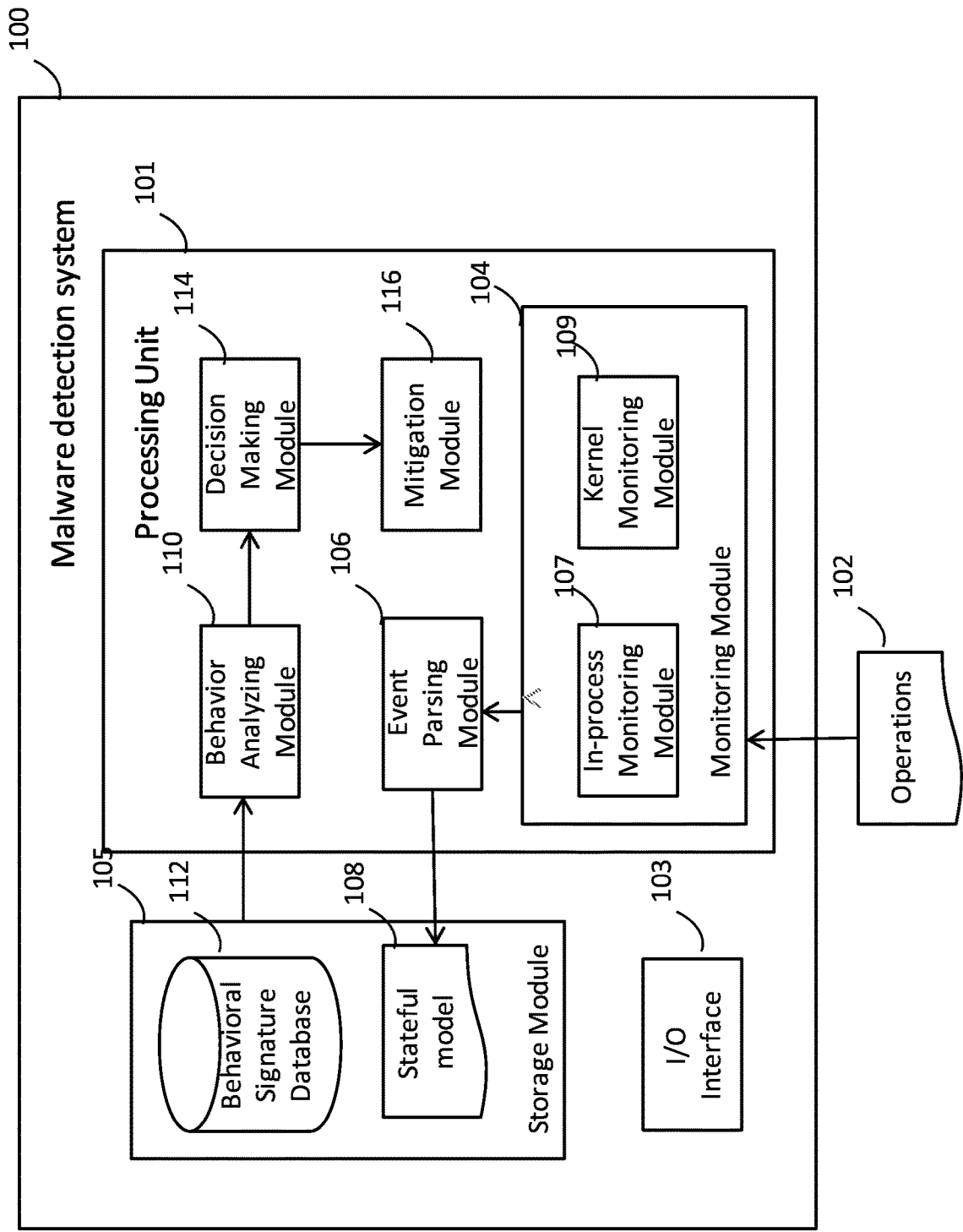


Fig. 1

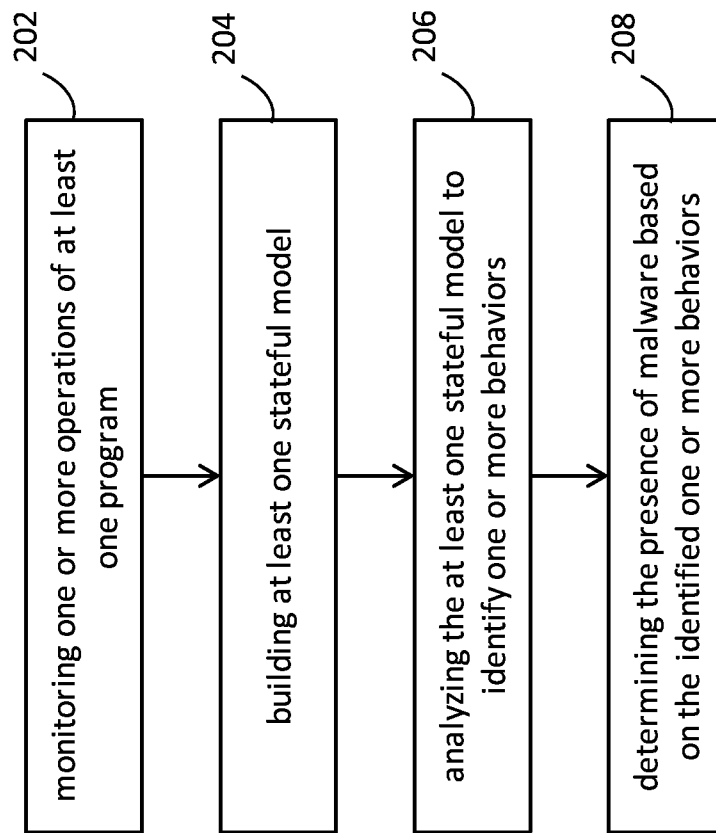


Fig. 2

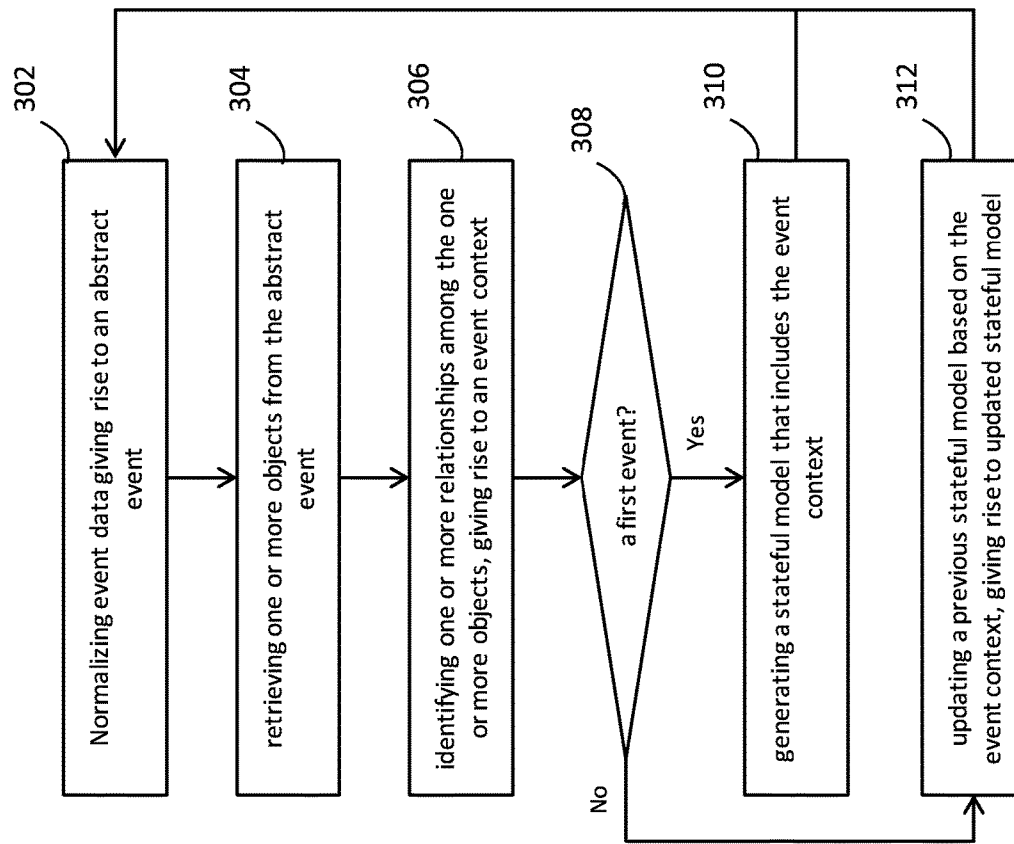


Fig. 3

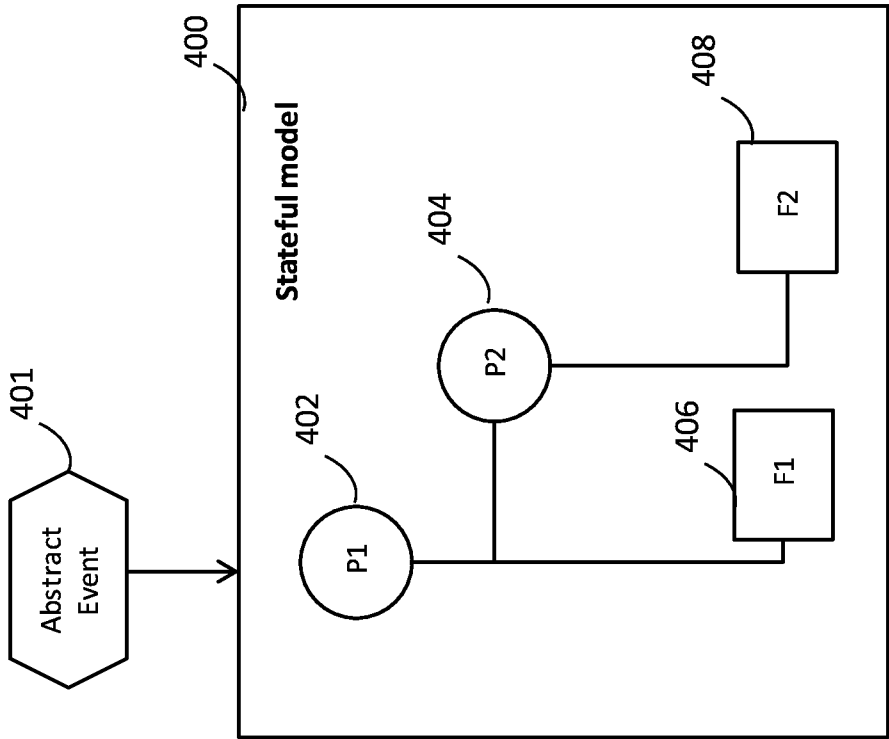


Fig. 4a

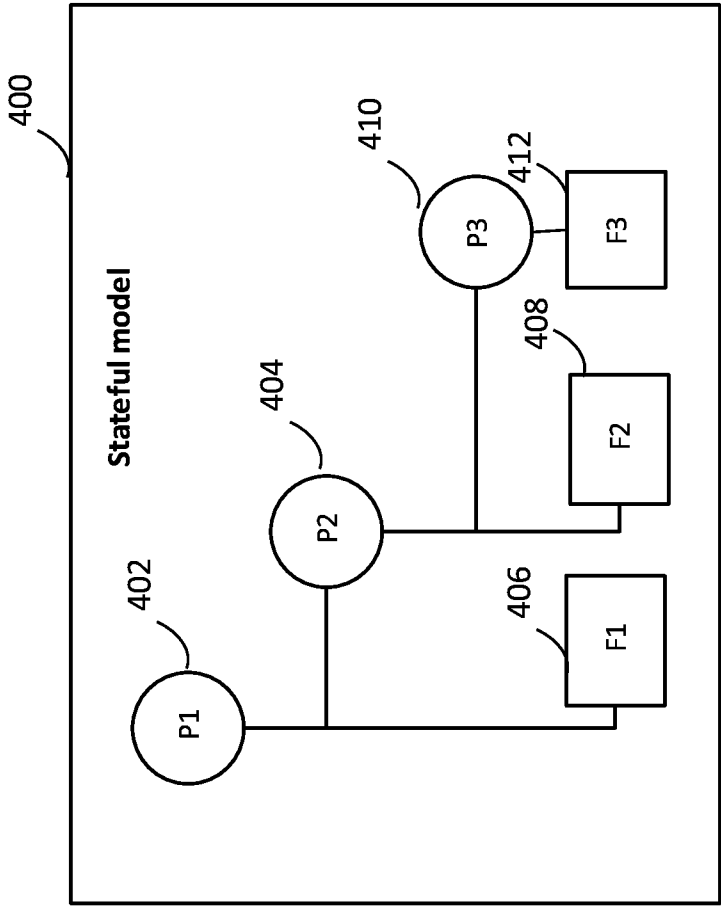


Fig. 4b

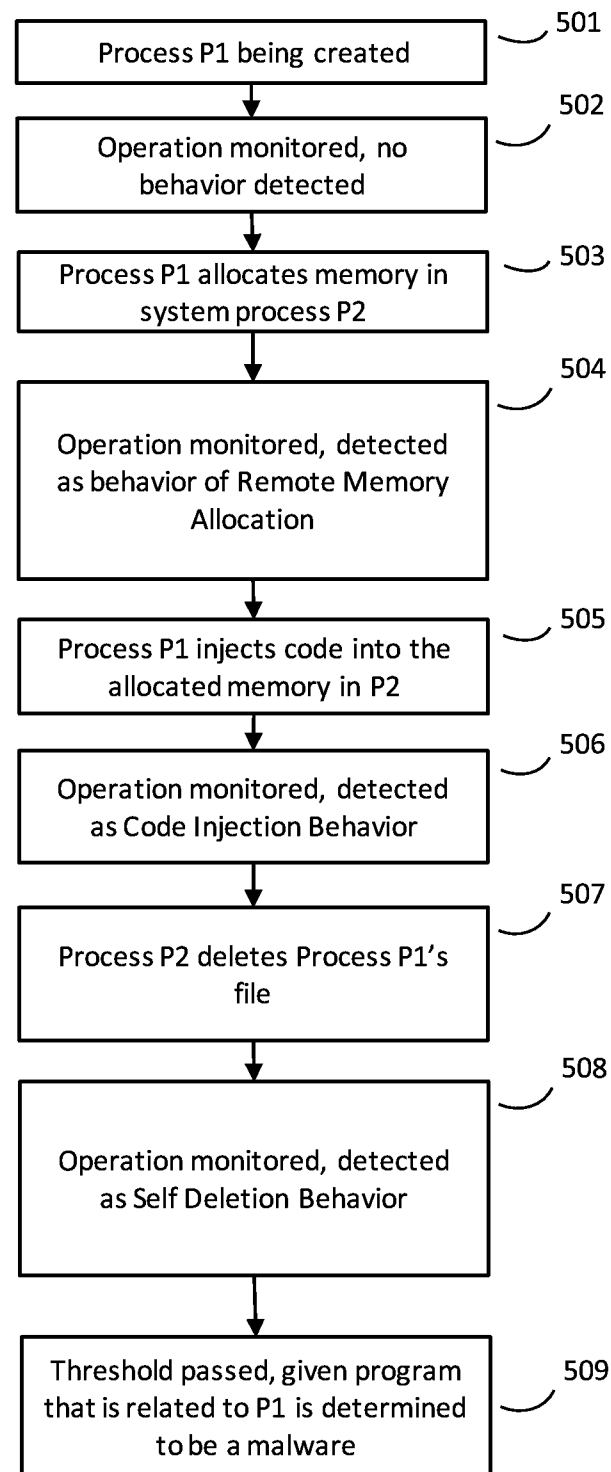


Fig. 5

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METHOD OF MALWARE DETECTION AND SYSTEM THEREOF

CROSS REFERENCE TO RELATED APPLICATIONS

The present application is a continuation of U.S. patent application Ser. No. 15/623,669, filed Jun. 15, 2017, now U.S. Pat. No. 10,664,596, issued May 26, 2020, which is a continuation of U.S. patent application Ser. No. 14/456,127, filed Aug. 11, 2014, now U.S. Pat. No. 9,710,648, issued Jul. 18, 2017, each of which is hereby incorporated herein by reference in its entirety under 37 C.F.R. § 1.57. Any and all applications for which a foreign or domestic priority claim is identified in the Application Data Sheet as filed with the present application are hereby incorporated by reference under 37 C.F.R. § 1.57.

TECHNICAL FIELD

The presently disclosed subject matter relates, in general, to the field of the detection of malicious software (malware), and more specifically, to methods and systems for behavior based malware detection.

BACKGROUND OF THE INVENTION

With the rapid growth of computer technology and widespread Internet access, malware threats have continued to grow significantly in recent decades, and thus have caused severe damage to systems, such as hardware failures and loss of critical data, etc.

Various antivirus technologies are currently in use, including signature and behavior based analysis, which aim to identify and prevent further spread of malware in the network. Signature-based analysis involves searching for known patterns of malicious code within executable code. However, malware is often modified (e.g., by obfuscating and randomizing content) in order to change its signature without affecting functionality, which renders the signature-based analysis mechanism as being increasingly ineffective. Due to an increase in malware variants (e.g., malware variants with the same behavior but different signatures), behavior-based analysis may be used to identify malware variants that have similar effects and thus can be handled with similar security measures.

Behavior-based analysis detects malware by monitoring behaviors of malicious activities rather than static signatures. Existing behavioral monitoring systems include a database of actions that are blacklisted and indicate malicious intent. If a given process or program performs any of the actions listed in the database, the action is blocked, and the process may be identified as malicious, and thus be terminated, by the monitoring system.

There is a need in the art for a new method and system for malware detection.

References considered to be relevant as background to the presently disclosed subject matter are listed below. Acknowledgement of the references herein is not to be inferred as meaning that these are in any way relevant to the patentability of the presently disclosed subject matter.

U.S. Pat. No. 8,555,385 (Bhatkar et al.) entitled “Techniques for behavior based malware analysis” discloses techniques for behavior based malware analysis. In one particular embodiment, the techniques may be realized as a method for behavior based analysis comprising receiving trace data, analyzing, using at least one computer processor, observable

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events to identify low level actions, analyzing a plurality of low level actions to identify at least one high level behavior, and providing an output of the at least one high level behavior.

U.S. Pat. No. 7,530,106 (Zaitsev et al.) entitled “System and method for security rating of computer processes” discloses a system, method, and computer program product for secure rating of processes in an executable file for malware presence, comprising: (a) detecting an attempt to execute a file on a computer; (b) performing an initial risk assessment of the file; (c) starting a process from code in the file; (d) analyzing an initial risk pertaining to the process and assigning an initial security rating to the process; (e) monitoring the process for the suspicious activities; (f) updating the security rating of the process when the process attempts to perform the suspicious activity; (g) if the updated security rating exceeds a first threshold, notifying a user and continuing execution of the process; and (h) if the updated security rating exceeds a second threshold, blocking the action and terminating the process.

U.S. Pat. No. 8,607,340 (Wright) entitled “Host intrusion prevention system using software and user behavior analysis” discloses improved capabilities for threat detection using a behavioral-based host-intrusion prevention method and system for monitoring a user interaction with a computer, software application, operating system, graphic user interface, or some other component or client of a computer network, and performing an action to protect the computer network based at least in part on the user interaction and a computer code process executing during or in association with a computer usage session.

US Patent Application No. 2012/079,596 (Thomas et al.) entitled “Method and system for automatic detection and analysis of malware” discloses a method of detecting malicious software (malware) including receiving a file and storing a memory baseline for a system. The method also includes copying the file to the system, executing the file on the system, terminating operation of the system, and storing a post-execution memory map. The method further includes analyzing the memory baseline and the post-execution memory map and determining that the file includes malware.

SUMMARY OF THE INVENTION

In accordance with an aspect of the presently disclosed subject matter, there is provided a computer-implemented method of detecting malware in real time in a live environment, the method comprising: monitoring one or more operations of at least one program concurrently running in the live environment; building at least one stateful model in accordance with the one or more operations; analyzing the at least one stateful model to identify one or more behaviors; and determining the presence of malware based on the identified one or more behaviors.

In accordance with an embodiment of the presently disclosed subject matter, there is further provided a computer-implemented method, wherein the monitoring further comprises: generating event data characterizing one or more events, each of the events being indicative of a corresponding monitored operation of the one or more operations.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a computer-implemented method, wherein the monitoring the one or more operations further includes selecting at least one operation of interest from the one or more operations, and monitoring the selected at least one operation of interest.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a computer-implemented method, wherein the at least one operation of interest includes one or more in-process operations and/or one or more kernel related operations.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a computer-implemented method, wherein the kernel related operations include one or more of the following: file system operations, process and memory operations, registry operations, and network operations.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a computer-implemented method, wherein the in-process operations are monitored by intercepting one or more library calls representing the in-process operations.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a computer-implemented method, wherein the kernel related operations are monitored by intercepting one or more system calls representing the kernel related operations.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a computer-implemented method, wherein the kernel related operations are monitored by registering one or more kernel filter drivers for the kernel related operations via one or more callbacks functions.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a computer-implemented method further comprising, for each event of the one or more events, generating a respective event data characterizing the event, wherein the event data includes at least the following attributes of the event: operation type, and source of the event.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a computer-implemented method, wherein the at least one stateful model includes one or more objects derived from the one or more operations and one or more relationships identified among the objects in accordance with the operations.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a computer-implemented method, wherein each of the objects represents an entity involved in the operations and is of a type selected from a group that includes: process object, file object, network object, registry object, and windows object.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a computer-implemented method, wherein each of the at least one stateful model is a program-level stateful model that represents a sequence of linked operations related to a given program of the at least one program.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a computer-implemented method, wherein the at least one stateful model is a system-level stateful model that represents operations related to all programs that run concurrently in the live environment.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a computer-implemented method, wherein the system-level stateful model includes one or more program-level stateful models each representing a sequence of linked operations related to a given program of the all programs.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a computer-implemented method further comprising monitoring

one or more kernel related operations of the at least one program; building at least one stateful model based on the monitored kernel related operations; analyzing the at least one stateful model to identify one or more behaviors; and determining the presence of malware based on a behavioral score of the stateful model.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a computer-implemented method wherein the at least one stateful model includes one or more objects derived from the one or more operations and one or more relationships identified among the objects in accordance with the operations.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a computer-implemented method, wherein the building the at least one stateful model comprises, for each event data associated with the event, normalizing the event data giving rise to an abstract event; retrieving one or more objects from the abstract event, each of the objects representing an entity involved in a corresponding operation and being of a type selected from a group that includes: process object, file object, network object, registry object and windows object, at least one of the objects representing the source of the event; identifying one or more relationships among the objects in accordance with the abstract event, and generating respective associations among the objects corresponding to the identified relationships, giving rise to an event context comprising the one or more objects and the associations therein; in case of the event being a first event of a stateful model, generating the stateful model including the event context; otherwise updating a previous stateful model based on the event context, giving rise to an updated stateful model, the previous stateful model corresponding to at least one previous event that precedes the event.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a computer-implemented method, wherein the updating further includes: in case the previous stateful model includes the one or more objects, adding the associations of the event context to the previous stateful model, giving rise to the updated stateful model; otherwise in case of at least one object of the objects being a new object that is not included in the previous stateful model, adding the new object and the associations of the event context to the previous stateful model, giving rise to the updated stateful model.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a computer-implemented method further comprising selecting selected event data associated with events of interest from the event data based on one or more predefined filtering rules and applying the normalizing of the event data with respect to the selected event data.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a computer-implemented method, wherein the one or more predefined filtering rules include filtering out event data associated with the following events: uncompleted events, memory related events of which a targeting process is not a remote process, and events in which a targeting process does not exist.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a computer-implemented method, wherein normalizing the event data includes formatting the event data and parsing the formatted event data giving rise to the abstract event.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a com-

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puter-implemented method, wherein the analyzing the at least one stateful model includes analyzing the event context in view of the stateful model or the updated stateful model in accordance with one or more predefined behavioral logics, and determining the presence of at least one behavior of the one or more behaviors upon any of the predefined behavioral logics being met, the at least one behavior related to a sequence of events of the stateful model including at least the event.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a computer-implemented method, wherein the predefined behavioral logics include determining a behavior of self-execution when the following condition is met: a target of an event is an object that is included in the stateful model.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a computer-implemented method wherein each of the at least one behavior is assigned with a respective behavioral score.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a computer-implemented method, wherein the determining the presence of malware further includes: in case of the at least one behavior being determined: searching if there is a previous stateful model score associated with the previous stateful model, the previous stateful model score being an aggregated behavioral score of all previous behavioral scores assigned for respective previous determined behaviors, the previous determined behaviors being related to the at least one previous event of the previous stateful model, if not, determining a sum of the respective behavioral score assigned for each of the at least one behavior as the stateful model score associated with the stateful model; otherwise increasing the previous stateful model score with the sum, giving rise to the stateful model score; comparing the stateful model score with a predefined threshold; and determining the presence of malware if the stateful model score passes the predefined threshold.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a computer-implemented method, wherein the respective behavioral score is assigned with a corresponding weight factor if a condition is met, and the increasing comprises applying the corresponding weight factor to the respective behavioral score giving rise to a respective weighted behavioral score, and increasing the previous stateful model score with a sum of the respective weighted behavioral score assigned for each of the at least one behavior.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a computer-implemented method, wherein the condition includes that a source of an event is a remote process and a target of the event is a system process.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a computer-implemented method, wherein the method further comprises: eliminating determined malware by remediating the one or more operations indicated by the stateful model.

In accordance with an aspect of the presently disclosed subject matter, there is yet further provided a system for detecting malware in real time in a live environment, the system comprising a processor configured to perform at least the following: monitor one or more operations of at least one program concurrently running in the live environment; build at least one stateful model in accordance with the one or more operations; analyze the at least one stateful model to

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identify one or more behaviors; and determine the presence of malware based on the identified one or more behaviors.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a system, wherein the at least one stateful model includes one or more objects derived from the one or more operations and one or more relationships identified among the objects in accordance with the operations.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a system, wherein each of the objects represents an entity involved in the operations and is of a type selected from a group that includes: process object, file object, network object, registry object, and windows object.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a system, wherein each of the at least one stateful model is a program-level stateful model that represents a sequence of linked operations related to a given program of the at least one program.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a system, wherein the at least one stateful model is a system-level stateful model that represents operations related to all programs that run concurrently in the live environment.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a system, wherein the system-level stateful model includes one or more program-level stateful models each representing a sequence of linked operations related to a given program of the all programs.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a system, wherein the processor is further configured to perform the following: monitor one or more kernel related operations of said at least one program; build at least one stateful model based on said monitored kernel related operations; analyze the at least one stateful model to identify one or more behaviors; and determine the presence of malware based on a behavioral score of said stateful model.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a system, wherein the kernel related operations include one or more of the following: file system operations, process and memory operations, registry operations, and network operations.

In accordance with an embodiment of the presently disclosed subject matter, there is yet further provided a system, wherein the processor is further configured to monitor the one or more kernel related operations by registering one or more kernel filter drivers for the kernel related operations via one or more callback functions.

In accordance with an aspect of the presently disclosed subject matter, there is yet further provided a non-transitory program storage device readable by machine, tangibly embodying a program of instructions executable by the machine to perform method steps for detecting malware in real time in a live environment, the method comprising: monitoring one or more operations of at least one program concurrently running in the live environment; building at least one stateful model in accordance with the one or more operations; analyzing the at least one stateful model to identify one or more behaviors; and determining the presence of malware based on the identified one or more behaviors.

BRIEF DESCRIPTION OF THE DRAWINGS

In order to understand the presently disclosed subject matter and to see how it may be carried out in practice, the

subject matter will now be described, by way of non-limiting example only, with reference to the accompanying drawings, in which:

FIG. 1 is a functional block diagram schematically illustrating a malware detection system, in accordance with certain embodiments of the presently disclosed subject matter;

FIG. 2 is a generalized flowchart of detecting malware in accordance with certain embodiments of the presently disclosed subject matter;

FIG. 3 is a generalized flowchart of building a stateful model in accordance with certain embodiments of the presently disclosed subject matter;

FIGS. 4a and 4b are schematic illustrations of an exemplified stateful model and an exemplified updated stateful model in accordance with certain embodiments of the presently disclosed subject matter, and

FIG. 5 is a generalized flowchart of an exemplified sequence of operations being monitored and processed in accordance with certain embodiments of the presently disclosed subject matter.

DETAILED DESCRIPTION

In the following detailed description, numerous specific details are set forth in order to provide a thorough understanding of the disclosed subject matter. However, it will be understood by those skilled in the art that the present disclosed subject matter can be practiced without these specific details. In other instances, well-known methods, procedures, and components have not been described in detail so as not to obscure the present disclosed subject matter.

In the drawings and descriptions set forth, identical reference numerals indicate those components that are common to different embodiments or configurations.

Unless specifically stated otherwise, as apparent from the following discussions, it is appreciated that throughout the specification discussions utilizing terms such as “monitoring”, “building”, “analyzing”, “determining”, “generating”, “selecting”, “normalizing”, “comparing”, “formatting”, “parsing”, “searching”, “increasing”, “eliminating”, “terminating”, “providing”, or the like, include action and/or processes of a computer that manipulate and/or transform data into other data, said data represented as physical quantities, e.g. such as electronic quantities, and/or said data representing the physical objects. The terms “computer”, “processor”, “processing unit”, “host machine”, and “end user station” should be expansively construed to include any kind of electronic device with data processing capabilities, including, by way of non-limiting examples, a personal computer, a server, a computing system, a communication device, a processor (e.g. digital signal processor (DSP), a microcontroller, a field programmable gate array (FPGA), an application specific integrated circuit (ASIC), etc.), any other electronic computing device, and or any combination thereof.

The operations in accordance with the teachings herein can be performed by a computer specially constructed for the desired purposes or by a general purpose computer specially configured for the desired purpose by a computer program stored in a non-transitory computer readable storage medium.

The term “non-transitory” is used herein to exclude transitory, propagating signals, but to otherwise include any volatile or non-volatile computer memory technology suitable to the presently disclosed subject matter.

As used herein, the phrase “for example,” “such as”, “for instance” and variants thereof describe non-limiting embodiments of the presently disclosed subject matter. Reference in the specification to “one case”, “some cases”, “other cases” or variants thereof means that a particular feature, structure or characteristic described in connection with the embodiment(s) is included in at least one embodiment of the presently disclosed subject matter. Thus the appearance of the phrase “one case”, “some cases”, “other cases” or variants thereof does not necessarily refer to the same embodiment(s).

It is appreciated that, unless specifically stated otherwise, certain features of the presently disclosed subject matter, which are described in the context of separate embodiments, can also be provided in combination in a single embodiment. Conversely, various features of the presently disclosed subject matter, which are described in the context of a single embodiment, can also be provided separately or in any suitable sub-combination.

In embodiments of the presently disclosed subject matter, fewer, more and/or different stages than those shown in FIGS. 2 and 3 may be executed. In embodiments of the presently disclosed subject matter one or more stages illustrated in FIGS. 2 and 3 may be executed in a different order and/or one or more groups of stages may be executed simultaneously. FIG. 1 illustrates a general schematic of the system architecture in accordance with an embodiment of the presently disclosed subject matter. Each module in FIG. 1 can be made up of any combination of software, hardware and/or firmware that performs the functions as defined and explained herein. The modules in FIG. 1 may be centralized in one location or dispersed over more than one location. In other embodiments of the presently disclosed subject matter, the system may comprise fewer, more, and/or different modules than those shown in FIG. 1.

As aforementioned, behavior-based analysis detects malware by monitoring behaviors of malicious activities rather than static signatures. There are a number of problems existing in current behavior-based technologies. For instance, due to the frequently changing behaviors of malicious programs, new instances of malwares may not be detected immediately due to lack of information about their behaviors and functionality. Current behavior-based technologies may also fail to trace a sequence of events, each of which, independently, is not identified as malicious, but when considered within the sequence context, is actually performing a malicious action. Moreover, current behavior-based technologies are normally implemented by performing emulation and running suspected malware in a safe environment (e.g., a sandboxed virtual machine) to reveal otherwise obscured logics and behaviors. This kind of emulation is normally very limited and the suspected malware under scrutiny never actually runs in a live environment. Thus it is impossible to actually observe full execution and interaction of the suspected malware with other processes and files that are not emulated in the safe environment. Therefore, not all potential malicious behaviors of the suspected malware can be detected by performing such emulation. Furthermore, it is typically resource-intensive to collect and analyze the large amount of operation information contained by suspicious malwares in order to identify potential behaviors, especially for a host machine with limited resources, such as an end user station. Certain embodiments of the detailed description are able to cope with these problems.

Bearing this in mind, attention is drawn to FIG. 1, schematically illustrating a functional block diagram of a

malware detection system in accordance with certain embodiments of the presently disclosed subject matter.

The term “malware” used in this specification should be expansively construed to include any kind of computer virus, ransomware, worms, trojan horses, rootkits, keyloggers, dialers, spyware, adware, malicious Browser Helper Objects (BHOs), rogue security software, or any other malicious or undesirable programs.

A Malware Detection System **100** illustrated in FIG. **1** implements a computer-based malware detection mechanism, which enables end users to detect malware in real time in a live environment. The term “live environment” used in this specification should be expansively construed to include any kind of host machine configuration where computer programs and products are actually put into operation for their intended uses by end users, such as, for example, an end user station with programs concurrently running in a production environment, in contrast to a safe environment, such as, for example, an emulated environment, or a sandboxed virtual machine environment.

As shown, the Malware Detection System **100** includes at least one Processing Unit **101** that comprises the following functional modules: Monitoring Module **104**, Event Parsing Module **106**, Behavior Analyzing Module **110**, and Decision Making Module **114**. Alternatively the Processing Unit **101** can be operatively coupled to the functional modules, and configured to receive instructions therefrom and execute operations in accordance with the instructions.

The Monitoring Module **104** can be configured to monitor, in real time, one or more operations **102** of at least one computer program that runs concurrently in the live environment. It is to be noted that the term “operation” used in this specification should be expansively construed to include any kinds of actions performed by one or more processes, threads, applications, files or any other suitable entities in any operating system. By way of non-limiting example, in a Windows operating system, operations can be performed by one or more processes of the computer programs. For purpose of illustration only, references are made in part of the following description with respect to operations performed by one or more processes. Embodiments are, likewise, applicable to operations performed by any other suitable entities in any operating system as described above.

A process is an instance of a computer program that is being executed. A process can further create child processes, and a computer program can be associated with one or more processes. It should be noted that the term “program” used in this specification should be expansively construed to include any kind of system software (e.g., operating system, device drivers, etc.) and application software (e.g., office suites, media players, etc.) that perform specified tasks with a computer.

As aforementioned, Monitoring Module **104** can monitor all the operations (e.g., performed by processes or other entities) occurred in the live system environment. According to certain embodiments, the Monitoring Module **104** can further include two sub-components: an In-process Monitoring Module **107** and a Kernel Monitoring Module **109**. The In-process Monitoring Module can monitor all in-process operations that are performed at process level and do not necessarily involve the kernel of an operating system. The Kernel Monitoring Module can monitor all operations that request services from an operating system’s kernel, such as file system operations, process and memory operations, registry operations, and network operations, as further elaborated with respect to FIG. **2**.

It is to be further noted that, without limiting the scope of the disclosure in any way, in some cases one operation can be construed to include a single action, such as “file read”. In some other cases, one operation can also be construed to include a sequence of actions, for example, “file copy” can be regarded as one operation which includes a sequence of three sequential actions “file create”, “file read”, and “file write”.

Event Parsing Module **106** can be configured to build at least one stateful model **108** in accordance with the one or more operations that are monitored by the Monitoring Module **104**. According to certain embodiments, a stateful model is a data model with hierarchical structure that contains information indicative of a real time updated system state resulted from a sequence of operations performed in a live environment. The sequence of operations can be linked together by context. Thus the stateful model can be a logical representation (e.g., a tree structure, etc) of a sequence of linked operations. For instance, the stateful model **108** can include one or more objects derived from real time operations **102**, and one or more relationships identified among the objects in accordance with the operations. According to certain embodiments, each of the objects of the stateful model **108** can represent an entity related in the operations and can be of a type selected from a group that includes: process object, file object, network object, registry object and windows object. The stateful model can further include attributes characterizing the objects and the identified relationships therein, as further elaborated with respect to FIGS. **3** and **4**.

Behavior Analyzing Module **110** can be configured to analyze the stateful model **108** constructed by Event Parsing Module **106** to identify one or more behaviors. It should be noted that the term “behavior” used in this specification should be expansively construed to include any sequence of operations performed by one or more processes that fulfill one or more predefined behavioral logics (also termed as “behavioral signatures” hereinafter).

According to certain embodiments, the Malware Detection System **100** can further comprise a Storage Module **105** that comprises a non-transitory computer readable storage medium. The Storage Module **105** can include a Behavioral Signature Database **112** that is operatively coupled to the Behavior Analyzing Module **110** and stores the one or more predefined behavioral logics. According to certain embodiments, the predefined behavioral logics are behavioral signatures indicative of specific behavioral patterns. In some cases, the behavioral logics can be predefined based on prior knowledge of certain malware behaviors, such as, for instance, self-deletion, self-execution, and code injection, etc. Optionally, the predefined behavioral logics can also include one or more logics indicative of benign behaviors, as further elaborated with respect to FIG. **2**. The stateful model **108** that is built by the Event Parsing Module **106** can also be stored in the Storage Module **105**.

Decision Making Module **114** can be configured to determine the presence of malware based on the one or more behaviors identified by the Behavior Analyzing Module **110**, as further elaborated with respect to FIG. **2**.

According to certain embodiments, the Processing Unit **101** can further include a Mitigation Module **116** configured to eliminate the determined malware by remediating the one or more operations indicated by the stateful model.

According to further embodiments, the Malware Detection System **100** can further include an I/O interface **103** communicatively coupled to the Processing Unit **101**. The I/O interface **103** can be configured to perform the following

actions: receive instructions from end users and/or from one or more of the functional modules, and provide an output of processed information obtained from the functional modules, e.g., an illustration of the determined malware, to the end users.

According to certain embodiments, the Processing Unit **101** is further configured to perform at least one of the aforementioned operations of the functional components of the Malware Detection System **100** in real time.

The operation of the Malware Detection System **100** and of the various components thereof is further detailed with reference to FIG. 2.

While not necessarily so, the process of operation of the Malware Detection System **100** can correspond to some or all of the stages of the method described with respect to FIG. 2. Likewise, the method described with respect to FIG. 2 and its possible implementations can be implemented by the Malware Detection System **100**. It is therefore noted that embodiments discussed in relation to the method described with respect to FIG. 2 can also be implemented, mutatis mutandis as various embodiments of the Malware Detection System **100**, and vice versa.

It should be further noted that the aforementioned functional components of the Malware Detection System **100** can be implemented in a stand-alone computer, such as the end user station. Or alternatively, one or more of the functional components can be distributed over several computers in different locations. In addition, the above referred modules can, in some cases, be cloud based.

Those versed in the art will readily appreciate that the teachings of the presently disclosed subject matter are not bound by the system illustrated in FIG. 1. Alternative to the example shown in FIG. 1, the Malware Detection System **100** can, in some cases, include fewer, more and/or different modules than shown in FIG. 1. Equivalent and/or modified functionality can be consolidated or divided in another manner and can be implemented in any appropriate combination of software, firmware and hardware.

Turning now to FIG. 2, there is shown a generalized flowchart of detecting malware in accordance with certain embodiments of the presently disclosed subject matter.

As illustrated in FIG. 2, one or more operations of at least one program that runs in a live environment can be monitored (**202**) in real time, e.g., by the Monitoring Module **104** of the Malware Detection System **100**. As aforementioned, in contrast to a safe environment, a live environment should include one or more computer programs that are put into operation for their intended uses. The computer programs run concurrently and interactively (e.g., with other programs and/or end users) in the live environment. According to certain embodiments, one or more processes can be launched by the one or more programs. Each process can perform one or more operations in order to communicate with and/or request services from the operating system. Accordingly, the Monitoring Module **104** can be configured to monitor the one or more operations performed by each process.

Due to the large number of concurrently running programs and operations thereof in a live environment, the amount of information contained in the monitored operations can be huge. According to certain embodiments, the Monitoring Module **104** can be configured to select at least one operation of interest from the one or more operations, and monitor the selected at least one operation of interest.

According to certain embodiments, the at least one operation of interest includes one or more in-process operations and/or one or more kernel related operations. In-process

operations can include any operation performed in user space (i.e., the memory area where application software executes) and do not necessarily involve the kernel of an operating system, such as, by way of non-limiting example, local process memory allocation, mapping functions from imported libraries, and read/write process memory. In some cases, the in-process operations can be monitored (e.g., by the In-process Monitoring module) by intercepting one or more library calls (e.g., API calls) that represent the corresponding operations. By way of non-limiting example, the In-process Monitoring module can attach monitoring hooks to the library calls in user space in order to monitor these calls.

The kernel related operations, on the other hand, can include one or more of the following operations that are performed in kernel space (i.e., the memory area reserved for running privileged kernel, kernel extensions, and most device drivers): file system operations, process and memory operations, registry operations, and network operations. Specifically, by way of non-limiting example, file system operations can include any operation and interaction with the storage medium of the host machine. Process and memory operations can include any operation of creating, terminating, modifying, querying, suspending and resuming processes, as well as memory management (e.g., allocating memory, creating a memory section, mapping/unmapping a memory section, writing/reading memory, etc). Registry operations can include any operation related to registry manipulation. And network operations can include any operation of sending or receiving data through network and network connection management.

The kernel related operations can be monitored by the Kernel Monitoring Module through different mechanisms, e.g., in accordance with different operating system platforms. For instance, for Mac OS X operating system, the kernel related operations can be monitored, by way of non-limiting example, by intercepting one or more system calls (in kernel space) that represent the corresponding operations. For the Windows operating system, kernel related operations can be monitored, by way of non-limiting example, by registering one or more kernel filter drivers for the kernel related operations via one or more callback functions. Windows operating system allows new drivers to be registered as part of the existing kernel stack, and thus information regarding a specific type of operation can be filtered by a corresponding kernel filter driver and passed through to the Kernel Monitoring Module via callback functions.

According to certain embodiments, OOB (Out-of-Band) monitoring approach can be adapted in the monitoring process (e.g., by the Kernel Monitoring Module). OOB enables the monitoring module to get notified on selected operations/events while not to have control over these operations/events, which allows the monitoring module to utilize different monitoring mechanisms (e.g., kernel callback functions) to accomplish a full system monitoring in an optimized manner. OOB also allows the events to be processed and analyzed into a stateful model in real time while the events are happening, as further described below.

According to certain embodiments, OOB can also enable the sequence of operations described with reference to FIG. 2, e.g., the monitoring operations, building stateful model, analyzing behaviors, determining malware and eliminating the determined malware, to be performed in the same machine, such as an end user station.

It is to be noted that the aforementioned categorized operations that are monitored respectively by different moni-

toring modules are provided for exemplary purposes only and should not be construed as limiting. For instance, in some cases one or more of the operations monitored by the In-process Monitoring Module can also be monitored by the Kernel Monitoring Module, and vice versa. According to one embodiment, at least one of the kernel related operations can be only monitored by the kernel Monitoring Module.

It should be noted that above mentioned examples of operations and implementations of the monitoring mechanisms are illustrated for exemplary purposes only. Additional kinds of operations and implementations can be applied in addition to or instead of the above.

It is also noted that the implementation mechanisms of the Kernel Monitoring Module can expedite system processing and enable the monitoring of the operations to be performed in a real time manner in a live environment.

According to certain embodiments, each monitored operation of the one or more operations constitutes an event. The Monitoring Module **104** can be further configured to generate event data characterizing one or more events. Optionally, an event data can be generated to characterize a respective event. According to certain embodiments, the event data can include at least the following attributes of the respective event: operation type, and source of the event.

Specifically, operation type is an identifier indicative of the type of the monitored operation that constitutes the event. The source of an event is the originating entity that performs the operation. Optionally, event data can include one or more additional attributes. For example, in some cases event data can include a target of an event, such as a targeting process, a targeting file, or any other entities that the operation is performed upon by the source of the event. In some further cases, event data can also include additional attributes according to different types of operations. For instance, event data that characterize file system operations can include additional attributes such as file permissions, full path of the file, size of the file, etc, while event data that characterize process and memory operations can include additional attributes such as address of the memory on which the operation is performed, size of the data that was written or read, memory permissions, etc.

Following step **202**, at least one stateful model can be built (**204**) in accordance with the one or more operations, e.g., by the Event Parsing Module **106** of the Malware Detection System **100**, as further described below in detail with respect to FIG. 3.

Attention is now directed to FIG. 3, illustrating a generalized flowchart of building a stateful model in accordance with certain embodiments of the presently disclosed subject matter. According to certain embodiments, the event data generated by the Monitoring Module **104** is created based on a large amount of raw data gathered through different routes, e.g., low level system calls and kernel driver callbacks, etc, thus the event data are generated in various forms. According to certain embodiments, this raw form of event data can be normalized (**302**) by the Event Parsing Module **106** into a logical data structure, giving rise to an abstract event which allows each segment of the attributes encoded in the event data to be accessed and analyzed. Specifically, the Event Parsing Module **106** can format the event data and parse the formatted event data in order to generate the abstract event. Through the event data normalization, event data indicative of similar operations but generated in various forms can also be normalized into a single format and categorized into the same event type. For example, various

system API calls generated to allocate memory will be categorized into a single type of abstract event, e.g., a memory allocation event.

According to certain embodiments, the Event Parsing Module **106** can select event data associated with events of interest from all event data received from the Monitoring Module **104** based on one or more predefined filtering rules, and apply the normalization with respect to the selected event data. By way of non-limiting example, the one or more predefined filtering rules can include filtering out event data associated with the following events: uncompleted events, memory related events in which the targeting process is not a remote process, and events in which the targeting process does not exist.

Based on the generated abstract event, a stateful model can be created or updated. As aforementioned, a stateful model can be a logical structure representation of a sequence of linked operations performed in a live environment.

According to certain embodiments, for each event data that is normalized to an abstract event, one or more objects can be retrieved (**304**) from the abstract event. As aforementioned, each of the retrieved objects represents an entity related in a corresponding event, and each object can be of a type selected from a group that includes: process object, file object, network object, registry object, and windows object. At least one of the objects represents the source of the event that performs a corresponding operation. By way of non-limiting example, the source of the event can be represented by a process object indicating an originating process that performs the operation. For example, a process **P1** performs an operation of "system shutdown". In this case, a process object will be retrieved from the corresponding abstract event to represent **P1** as the source of the event.

In some cases an operation is performed upon a target entity (i.e. target of the event) by the source of the event. For example, a process **P1** opens a file **F1**. A process object will be retrieved from the corresponding abstract event to represent **P1** as the source of the event for the operation "file open", and a file object will be retrieved to represent **F1** as the target of the event.

It is to be noted that an operation is usually initiated by a process. Thus the source of an event is normally represented by a process object. The target of the event, however, can be of various types of objects that are manipulated in the operation, such as a process object, file object, network object, registry object, etc.

According to further embodiments, a process can own resources, such as a source file that the process is initiated from. The source file can be of various types, such as, by way of non-limiting example, a document file, an image file that contains the executable program that the process is launched from, or any other relevant types of files. A source file, if related to an operation, can also be represented by a file object.

It is to be noted that the above mentioned object types are merely illustrated for exemplary purposes only and should not be construed as limiting the present disclosure in any way. Additional types of objects that may occur in an operation can be included in addition to or instead of the above.

Following retrieving the objects from an abstract event in step **304**, the Event Parsing Module **106** can identify one or more relationships (**306**) among the objects in accordance with the abstract event, and generate respective associations among the objects corresponding to the identified relationships, giving rise to an event context corresponding to the abstract event. The event context contains context informa-

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tion of the corresponding event, and comprises the one or more objects of the event and the associations therein.

Following step 306, the Event Parsing Module 106 can further determine if a current event is a first event (308) of a stateful model, as described below in detail with respect to FIG. 4a. In case of the above condition being met, a new stateful model can be generated (310) and include the event context, namely, the one or more objects and the one or more associations therein. The process then goes back to step 302 wherein the next event data can be processed.

With reference now to FIG. 4a, there is shown an exemplified stateful model 400 being created based on an abstract event 401, in accordance with certain embodiments of the presently disclosed subject matter. Abstract event 401 is normalized from an event data characterizing an event E1 of a process P1 creating a child process P2. The abstract event 401 comprises the following attributes of the event: operation type—process creation; source of the event—P1 (as the originating process of the event), source file of P1—F1, target of the event—P2 (as a targeting process of the event), and source file of P2—F2. Based on the abstract event 401, four objects can be retrieved: a process object 402 indicative of the source of the event P1, a process object 404 indicative of the target of the event P2, a file object 406 indicative of the source file F1 of P1, and a file object 408 indicative of the source file F2 of P2. According to certain embodiments, file objects 406 and 408 can be affiliated with, or correlated with, their respective process objects 402 and 404 as illustrated. The abstract data 401 can further include additional attributes which contain more information of the operation if applicable.

A relationship indicative of process creation can be identified between process objects 402 and 404 in accordance with the abstract event. A corresponding association between 402 and 404 can be generated accordingly based on the identified relationship, giving rise to an event context that comprises the process objects 402 and 404 (together with their correlated file objects 406 and 408) and the association therebetween. The association can be represented, e.g., as a direct linkage between the two related objects 402 and 404, as illustrated in FIG. 4a.

According to certain embodiments, one or more fields can be created for each of the objects, storing one or more parameters characterizing the respective object and the association related thereto. By way of non-limiting example, the process object 402 can have one or more fields selected from a group that includes: process identifier (e.g., a unique identifier assigned by the operating system for each process), one or more source file identifiers (e.g., a pointer to file object 406), and one or more operations and corresponding associations related thereto (e.g., an operation of process creation and a corresponding linkage to P2). The file object 406 can have one or more of fields selected from a group that includes: file identifier (e.g., the full path of the file), process identifier, and one or more operations and corresponding associations related thereto. Assume that E1 is a first event in a stateful model, a stateful model 400 can be generated and include the event context of E1.

It should be noted that the term “stateful model” should be expansively construed to include any of the following situations:

1) A stateful model can be a program-level stateful model that represents a sequence of linked operations related to a given program (and in some cases, also operations related to one or more other programs that are linked to the given program due to operations). In this case, a stateful model

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represents a program context that reflects all the operations related to the given program by context.

A first event of the program-level stateful model can be determined to be any event that relates to the given program’s first interaction with the system. For instance, a first event can be determined to be an event of “process creation” that creates the initiating process of the given program. An initiating process is the process that is created upon the given program being executed, which may also be the root process of a stateful model that performs further operations. A first event can also be determined to be an event performed by the initiating process upon other objects.

In the above example illustrated in FIG. 4a, if the originating process P1 is the initiating process of a certain program, the creation of P1 can be determined as the first event in the stateful model. Since the initiating process may be created by a system process P0, in some cases the stateful model can include P0, P1 and the association of process creation between P0 and P1. In some other cases the stateful model may include only the object P1, and a reference therewith indicating that P0 is the parent of P1. In some further cases a first event can also be determined as an event that P1 performs on other objects, for example, an event of “process creation” performed by P1 to create a child process P2.

In some circumstances events can be delayed to be processed by the Event Parsing Module 106 due to unexpected system processing problems. Thus a first event of the stateful model can also be an event that does not occur first in terms of time, but is first processed by the Event Parsing Module 106. Accordingly, following the above mentioned example of FIG. 4a, if a further event E2 of P2 opening a file F1 is first processed by the Event Parsing Module 106, the event E2 can be determined to be a first event of the stateful model, and any event that occurs before it (e.g., the event E1 of P1 creating P2) can be processed retroactively and reflected in the stateful model.

Thus, depending on the number of programs concurrently running in the live environment and the operational relationships among them, there may be one or more program stateful models co-existing, each of which represents a respective program context of a given program;

2) A stateful model can be a system-level stateful model that represents operations related to all programs that run concurrently in a live environment. In this case a first event of the stateful model can be determined to be the event of “system start” that is initiated when the operating system initially starts. Accordingly, there is only one stateful model existing at any given time in the system which represents a system context of the entire environment. According to some embodiments, the system-level stateful model can be created upon the initialization of the operating system, and can be kept updating while the operating system and program processing proceeds. In accordance with further embodiments, the system-level stateful model may be created by including one or more program-level stateful models each related to one program of all the programs running in the live environment as described above.

It is to be noted that the definition and implementation of the above stateful model structure are illustrated for exemplary purposes only and should not be construed as limiting the present disclosure in any way. Alternative data structures can be applied to implement equivalent functionality of the stateful model in addition to or in lieu of the above.

Turning back to FIG. 3, according to certain embodiments, if the current event is not a first event of a stateful model (308), a previous stateful model corresponding to at

least one previous event that precedes the current event exists. The Event Parsing Module 106 can update (312) the previous stateful model based on the event context of the current event, giving rise to an updated stateful model that is updated.

According to certain embodiments, a previous stateful model can be updated in accordance with the following scenarios:

1) If all the objects of the current event are already included in the previous stateful model, the one or more associations of the event context can be added to the previous stateful model, giving rise to the updated stateful model;

2) Otherwise at least one object of the one or more objects should be a new object that does not exist in the previous stateful model. Thus the new object, together with the one or more associations, can be added to the previous stateful model, giving rise to the updated stateful model.

Continuing with the example illustrated in FIG. 4a, assume that the illustrated stateful model 400 (including process objects P1, P2 and the association between P1 and P2 representing the event E1 of P1 creating P2) is a previous stateful model that exists, and a current event E2 arrives, wherein the same process P1 allocates memory in the same child process P2. Following the process in FIG. 3, the event data that characterizes the current event E2 is normalized to an abstract event. Objects P1 and P2 are retrieved based on the abstract event. A relationship indicative of memory allocation can be identified between P1 and P2 based on the abstract event, and an association between P1 and P2 can be generated based on the identified relationship. Thus an event context for the current event E2 comprises objects P1 and P2 and the association therebetween. Since the current event E2 is not a first event in the previous stateful model 400, the stateful model 400 will be updated based on the current event context. In this case, since all the objects of the current event, namely, P1 and P2, are already included in the previous stateful model 400, the currently generated association between P1 and P2 representing an operation of memory allocation, will be added as a new association between P1 and P2 in the stateful model 400, besides the previous association therebetween representing the operation of process creation, giving rise to an updated stateful model. By way of non-limiting example, the new association can be added in the stateful model by adding a respective field for P1 and/or P2 to indicate the operation of memory allocation therebetween. Since only the association has been updated, the hierarchical structure of the updated stateful model may look similar as illustrated in FIG. 4a, with a newly added association.

Continuing with the same example, assume that another event E3 arrives after E2, wherein the process P2 creates a child process P3. Following the same process in FIG. 3, the event data that characterizes the current event E3 is normalized to an abstract event. Objects P2 and P3 are retrieved based on the abstract event. A relationship indicative of process creation can be identified between P2 and P3 based on the abstract event, and an association between P2 and P3 can be generated based on the identified relationship. Thus an event context for the current event E3 comprises objects P2 and P3 and the association therebetween. Since the current event E3 is not the first event in the stateful model 400, the stateful model 400 will be updated based on the current event context. In this case, since P3 is a new object that does not exist in the previous stateful model, the new object P3 can be added to the stateful model 400 as a process object 410. Optionally a file object F3 that is correlated with

P3 can also be added as a file object 412. The association between P2 and the new object P3 can be added in the stateful model, by way of non-limiting example, by adding a respective field for P2 and/or P3 to indicate the operation of process creation therebetween, giving rise to an updated stateful model, as illustrated in FIG. 4b.

It is to be noted that the specific examples of building and updating the stateful model illustrated above are provided for exemplary purposes only and should not be construed as limiting. Accordingly, other ways of implementation of building and updating the stateful model can be used in addition to or in lieu of the above.

It should also be noted that the present disclosure is not bound by the specific sequence of operation steps described with reference to FIG. 3.

Having described the structure of the stateful model and the process of building/updating the stateful model in accordance with certain embodiments, attention is now drawn back to FIG. 2, wherein analyzing at least one stateful model in order to identify one or more behaviors is now described with reference to step 206.

According to certain embodiments, the Behavior Analyzing Module 110 can be further configured to analyze the event context of the current event in view of the stateful model (when the stateful model is newly created based on the current event) or the updated stateful model (when the stateful model is updated based on the current event), in accordance with one or more predefined behavioral logics.

The Behavior Analyzing Module 110 can further determine the presence of at least one behavior upon any of the one or more predefined behavioral logics being met. The determined behavior relates to a sequence of events of the stateful model including at least the current event. In some cases, each of the sequence of events independently may not be identified as malicious, but when considered within the sequence context, is actually performing a malicious behavior. By analyzing the event context in view of the stateful model, the Behavior Analyzing Module can inspect a specific event while looking at the whole picture, thus avoiding omission of undetected malwares.

According to certain embodiments, the predefined behavioral logics are behavioral signatures indicative of specific behavioral patterns. The behavioral logics can be predefined based on prior knowledge of certain malware behaviors, such as, for instance, self-deletion, self-execution, and code injection, etc. The behavioral logics can be stored in a Behavioral Signature Database 112 as aforementioned with respect to FIG. 1. One of the predefined behavioral logics can be, by way of non-limiting example, determining a behavior of self-execution when the following condition is met: the target of an event is an object that is already included in the stateful model, which indicates that the operation is performed on one of the family's own members. For example, a targeting process of an event is identical to one of the following: the originating process, a child process of the originating process, and a parent process of the originating process. Another similar exemplary behavioral logic can be, for instance, determining a behavior of self-deletion when the following condition is met: the target of a deletion event is an object included in the stateful model. For example, the source file of a targeting process of an event is identical to one of the following: the source file of the originating process, and the source file of a parent process of the originating process.

Optionally, the predefined behavioral logics can also include one or more logics indicative of benign behavior patterns such as, for example, interaction with the desktop or

users, registration in the system program repository, etc. According to certain embodiments, each behavioral signature in the database can be associated with a predefined behavioral score that indicates the malicious level of a corresponding behavior. Accordingly each of the determined at least one behavior can be assigned with a respective behavioral score based on the predefined behavioral score associated therewith. The process of analyzing a stateful model and determining at least one behavior is further exemplified with reference to FIG. 5.

It is to be noted that the hierarchical structure of the stateful model as described above is designed as a fast accessible data structure, which can in turn enable the creating of the stateful model and analyzing the created stateful model, following the monitoring of the operations, to be performed in a real time manner in a live environment.

Upon the at least one behavior being determined, the Decision Making Module 114 can be configured to determine the presence of malware (208) based on the determined behavior. According to certain embodiments, each stateful model can be associated with a stateful model score. The stateful model score is an aggregated behavioral score of all behavioral scores assigned for respective behaviors being determined in the stateful model. Upon at least one current behavior being determined in a stateful model, the Decision Making Module 114 can search if there is a previous stateful model score associated with a previous stateful model. Accordingly, the previous stateful model score is an aggregated behavioral score of all previous behavioral scores assigned for respective previous determined behaviors, the previous determined behaviors being related to the at least one previous event of the previous stateful model. If there is no previous stateful model score, the sum of the respective behavioral score for each of the at least one behavior can be determined as the stateful model score associated with the current stateful model. Otherwise, if there is found a previous stateful model score, the previous stateful model score can be increased with the sum, giving rise to the stateful model score that has been updated based on the current event. The Decision Making Module 114 can be further configured to compare the stateful model score with a predefined threshold. The predefined threshold can be a score indicative of malware presence and can be predetermined based on prior knowledge of malware detection. If the stateful model score passes the predefined threshold, a presence of malware can be determined. For example, the corresponding stateful model, and one or more programs that relate to the stateful model can be determined as malicious. The process of determining the presence of malware is further exemplified with reference to FIG. 5.

According to certain embodiments, the respective behavioral score of a currently determined behavior can be assigned with a corresponding weight factor if a condition is met. The condition can be, by way of non-limiting example, that the source of an event is a remote process and the target of the event is a system process, indicating that a remote process is performing operations on a system process. In this case a weight factor (e.g., a numerical value greater than 1) can be assigned to the original behavioral score associated with this behavior, indicating an increasing likelihood of malware presence. The assigned weight factor can be applied to the original behavioral score (e.g., by multiplying the original behavioral score with the weight factor), giving rise to a weighted behavioral score. Accordingly the previous stateful model score can be increased with a sum of the weighted behavioral score assigned for each of the at least one behavior.

It is to be noted that the present disclosure is not bound by the specific scoring and weighting paradigms described above. The scoring and weighting functionalities can be implemented in a consolidated manner or separately. Additional kinds of implementations can be applied in addition or instead of the above.

According to certain embodiments, the determined malware can be eliminated by remediating the operations indicated in the stateful model, such as, by way of non-limiting example, by terminating the process objects within the stateful model, removing the file objects (or other types of objects) correlated with or created by the process objects, and undo the operations performed among the objects if possible, in order to restore the system to a state as close as possible to the system state before the operations in the stateful model were executed. According to further embodiments, an output of the determined malware can be provided through the I/O Interface 103 to the end users, as aforementioned.

According to certain embodiments, the sequence of operations described with reference to FIG. 2, e.g., the monitoring operations, building stateful model, analyzing behaviors, determining malware and eliminating the determined malware, can be carried out concurrently in real time. For instance, building at least one stateful model in accordance with the one or more operations responsive to monitoring the one or more operations of at least one program concurrently running in a live environment can be performed in real time. Additionally or alternatively, analyzing the at least one stateful model to identify one or more behaviors responsive to monitoring the one or more operations and building the at least one stateful model can be performed in real time. Additionally or alternatively, determining the presence of malware based on the identified one or more behaviors responsive to analyzing the at least one stateful model can be performed in real time. Additionally or alternatively, eliminating the determined malware responsive to determining the presence of malware can be performed in real time.

It is to be noted that the present disclosure is not bound by the specific sequence of operation steps described with reference to FIG. 2.

Turning now to FIG. 5, there is shown a generalized flowchart of an exemplified sequence of operations being monitored and processed in accordance with certain embodiments of the presently disclosed subject matter.

As shown, a process P1 is created (501) upon a given program being executed. Thus P1 is the initiating process of the given program. The operation of process creation is monitored, e.g., by the kernel monitoring module. A corresponding event E1 and event data thereof are generated accordingly. E1 is determined to be the first event of a stateful model, and the stateful model is generated based on E1. The stateful model will now include an event context of E1, namely, P1 (optionally, also a system process P0 that creates P1, and/or the source file F1 of P1), together with an association of process creation of P1. No behavior is determined (502) at this stage in accordance with the predefined behavioral logics, and, accordingly no score is assigned.

A second operation of P1 allocating memory to a system process P2 (503) occurs. The operation of memory allocation is monitored, e.g., by the in-process monitoring module. A corresponding event E2 and event data thereof are generated accordingly. Since E2 is not the first event of a stateful model, the previous stateful model comprising event context E1 is updated based on E2. The stateful model now includes P1, P2 (optionally also their source files F1 and F2) together

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with an association of memory allocation between P1 and P2. A behavior of remote memory allocation is determined (504) in accordance with one of the predefined behavioral logics, and accordingly a behavioral score S1 is assigned. Since there is no previous stateful model score, the behavioral score S1 is also the stateful model score.

Following the second operation, a third operation of P1 injecting code (505) in the allocated memory in P2 occurs. According to certain embodiments, the operation of code injection can comprise three actions: memory write, memory execution permissions, and code execution, all of which are monitored. A corresponding event E3 and event data thereof are generated accordingly. Since E3 is not the first event of a stateful model, the previous stateful model based on event context of E1 and E2 are further updated based on the current event E3. The stateful model now includes P1, P2 (optionally also their source files F1 and F2), a previous association of memory allocation between P1 and P2, and a new association of code injection between P1 and P2. A behavior of code injection is determined (506) in accordance with one of the predefined behavioral logics, and accordingly a behavioral score S2 is assigned. The stateful model score is updated to be the sum of S1 and S2.

A fourth operation of P2 deleting P1's file F1 (507) follows the third operation. The operation of file deletion is monitored. A corresponding event E4 and event data thereof are generated accordingly. Since E4 is not the first event of a stateful model, the previous stateful model based on previous events E1, E2 and E3 are now updated based on E4. The present stateful model includes P1, P2, F1 (optionally also source file F2), two associations (i.e. memory allocation, and code injection) between P1 and P2, and a new association of file deletion between P2 and F1. Based on analyzing the stateful model, it is noted that P1 is actually the parent of P2. A behavior of self-deletion is determined (508) in accordance with one of the predefined behavioral logics, and a behavioral score S3 is assigned. Now the stateful model score is updated to be the sum of S1, S2 and S3. If the stateful model score passes a predefined threshold, the presence of malware is determined. For example, the stateful model, especially the given program that is related to P1 is determined to be malicious, and will be eliminated (509). For instance, the process objects P1 and P2 are terminated, the file objects F1 and F2 are removed, and the relevant operations between P1 and P2, such as memory allocation, code injection, file deletion etc, can be remediated if possible.

It is to be noted that the specific examples illustrated above with reference to FIG. 5 are provided for exemplary purposes only and should not be construed as limiting the present disclosure in any way.

It is to be understood that the presently disclosed subject matter is not limited in its application to the details set forth in the description contained herein or illustrated in the drawings. The presently disclosed subject matter is capable of other embodiments and of being practiced and carried out in various ways. Hence, it is to be understood that the phraseology and terminology employed herein are for the purpose of description and should not be regarded as limiting. As such, those skilled in the art will appreciate that the conception upon which this disclosure is based can readily be utilized as a basis for designing other structures, methods, and systems for carrying out the several purposes of the present presently disclosed subject matter.

It will also be understood that the system according to the presently disclosed subject matter can be implemented, at least partly, as a suitably programmed computer. Likewise,

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the presently disclosed subject matter contemplates a computer program being readable by a computer for executing the disclosed method. The presently disclosed subject matter further contemplates a machine-readable memory tangibly embodying a program of instructions executable by the machine for executing the disclosed method.

The invention claimed is:

1. A computer-implemented method of performing a behavior-based analysis of an execution of a program in an operating system, the method comprising:

monitoring, by a computer system, by registering one or more kernel filter drivers for kernel space operations via one or more call back functions using an out-of-band monitoring module, one or more operations performed by the execution of the program running in the operating system in a live environment, wherein the monitoring comprises tracking user space operations and the kernel space operations;

selecting at least one operation of interest from the one or more operations;

generating, by the computer system, an event data for each of the at least one operation of interest, wherein the event data characterizes one or more events of the at least one operation of interest;

filtering event data of interest from the event data for each of the at least one operation of interest, the filtering based on one or more predefined filtering rules;

normalizing the event data of interest into a logical data structure such that attributes of the event data of interest can accessed and analyzed;

building, by the computer system, at least one stateful model of the execution of the program based on the normalized event data of interest, the at least one stateful model comprising a hierarchical structure of the at least one operation of interest performed by the execution of the program in the live environment, the at least one operation of interest linked by an event context, wherein the hierarchical structure comprises:

the event context comprising:

one or more objects derived from the one or more monitored operations;

one or more fields generated for each of the one or more objects, the one or more fields storing one or more parameters characterizing a respective object of the one or more objects and an associate to the respective object; and

one or more relationships identified among the one or more objects; and

attributes characterizing the one or more objects and the one or more relationships among the one or more objects, wherein the attributes comprise at least a type of the at least one operation of interest and a source of the one or more events, wherein the type comprises an identifier of the at least one operation of interest that characterizes the one or more events, and wherein the source comprises an originating entity that performs the at least one operation of interest,

wherein each of the one or more objects represent an entity related to the one or more monitored operations;

analyzing, by the computer system, the event context in view of the at least one stateful model to identify one or more behaviors of the execution of the program related to the one or more events;

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applying a score to the stateful model based on the one or more identified behaviors, wherein applying the score to the stateful model comprises:

- determining a behavior score for each of the one or more identified behaviors;
- assigning a weight factor to each behavior score associated with the one or more identified behaviors to generate a weighted behavior score for each of the one or more identified behaviors, wherein the weighted behavior score indicates the likelihood of the presence of malware based on the one or more identified behaviors;
- determining a sum of the weighted behavior scores for each of the one or more identified behaviors; and
- comparing the one or more identified behaviors and the score to one or more pre-existing behaviors and a pre-existing score of a pre-existing stateful model, wherein the computer system comprises a processor and memory.

2. The method of claim 1, further comprising updating, in real time, the at least one stateful model in response to one or more new events.

3. The method of claim 1, further comprising outputting, via an output device of the computer system, a representation of the one or more identified behaviors of the execution of the program.

4. The method of claim 1, further comprising storing the one or more identified behaviors of the execution of the program in a behavioral profile database.

5. The method of claim 1, wherein the computer system comprises a cloud-based computer system.

6. The method of claim 1, wherein the computer system comprises one or more functional components distributed over more than one computer.

7. The method of claim 1, wherein the live environment comprises one or more programs, including the program, operating concurrently and interactively for their intended uses.

8. The method of claim 1, further comprising aggregating the one or more identified behaviors.

9. The method of claim 1, wherein the one or more behaviors comprise a representation of a behavior pattern of the execution of the program.

10. The method of claim 1, further comprising analyzing the one or more behaviors to determine if the execution of the program comprises malware.

11. A system for performing a behavior-based analysis of an execution of a program in an operating system, the system comprising:

- one or more computer readable storage devices configured to store a plurality of computer executable instructions; and
- one or more hardware computer processors in communication with the one or more computer readable storage devices and configured to execute the plurality of computer executable instructions in order to cause the system to:
 - monitor, by registering one or more kernel filter drivers for kernel space operations via one or more call back functions using an out-of-band monitoring module, one or more operations performed by the execution of the program running in the operating system in a live environment, wherein monitoring comprises tracking user space operations and the kernel space operations;
 - select at least one operation of interest from the one or more operations;

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generate an event data for each of the at least one operation of interest, wherein the event data characterizes one or more events of the at least one operation of interest;

filter event data of interest from the event data for each of the at least one operation of interest, the filtering based on one or more predefined filtering rules;

normalize the event data of interest into a logical data structure such that attributes of the event data of interest can accessed and analyzed;

build at least one stateful model of the execution of the program based on the normalized event data of interest, the at least one stateful model comprising a hierarchical structure of the at least one operation of interest performed by the execution of the program in the live environment, the at least one operation of interest linked by an event context, wherein the at least one stateful model comprises:

- the event context comprising:
 - one or more objects derived from the one or more monitored operations;
 - one or more fields generated for each of the one or more objects, the one or more fields storing one or more parameters characterizing a respective object of the one or more objects and an associate to the respective object; and
 - one or more relationships identified among the one or more objects; and
- attributes characterizing the one or more objects and the one or more relationships among the one or more objects, wherein the attributes comprise at least a type of the at least one operation of interest and a source of the one or more events, wherein the type comprises an identifier of the at least one operation of interest that characterizes the one or more events, and wherein the source comprises an originating entity that performs the at least one operation of interest,

wherein each of the one or more objects represent an entity related to the one or more monitored operations;

analyze the event context in view of the at least one stateful model to identify one or more behaviors of the execution of the program related to the one or more events;

apply a score to the stateful model based on the one or more identified behaviors, wherein applying the score to the stateful model comprises:

- determining a behavior score for each of the one or more identified behaviors;
- assigning a weight factor to each behavior score associated with the one or more identified behaviors to generate a weighted behavior score for each of the one or more identified behaviors, wherein the weighted behavior score indicates the likelihood of the presence of malware based on the one or more identified behaviors;
- determining a sum of the weighted behavior scores for each of the one or more identified behaviors; and
- compare the one or more identified behaviors and the score to one or more pre-existing behaviors and a pre-existing score of a pre-existing stateful model.

12. The system of claim 11, wherein the system is further caused to update, in real time, the at least one stateful model in response to one or more new events.

13. The system of claim 11, wherein the system is further caused to output, via an output device of the system, a representation of the one of more behaviors of the execution of the program.

14. The system of claim 11, wherein the system is further 5 caused to store the one or more behaviors of the execution of the program in a behavioral profile database.

15. The system of claim 11, wherein the system comprises a cloud-based computer system.

16. The system of claim 11, wherein the system comprises 10 one or more functional components distributed over more than one computer.

17. The system of claim 11, wherein the live environment comprises one or more programs, including the program, operating for concurrently and interactively for their 15 intended uses.

18. The system of claim 11, wherein the system is further caused to aggregate the one or more identified behaviors.

19. The system of claim 11, wherein the one or more behaviors comprise a representation of a behavior pattern of 20 the execution of the program.

20. The system of claim 11, wherein the system is further caused to analyze the one or more behaviors to determine if the execution of the program comprises malware.

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